

# Decidable $\mathbf{0}'$ -categoricity of models which realize only types with low CB ranks

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Logic Colloquium 2018

# Decidable categoricity

## Definition

A computable structure  $\mathcal{A}$  is **d-computably categorical** (**d-autostable**) if for every computable structure  $\mathcal{B}$  isomorphic to  $\mathcal{A}$ , there exists a d-computable isomorphism from  $\mathcal{A}$  onto  $\mathcal{B}$ .

Goncharov investigated computable categoricity, restricted to decidable structures.

## Definition

A decidable structure  $\mathcal{A}$  is called **decidably d-categorical** (**d-autostable relative to strong constructivizations**) if every two decidable copies of  $\mathcal{A}$  are d-computably isomorphic.

# Categoricity spectrum

Definition( Fokina, Kalimullin, Miller, 2010)

The **categoricity spectrum** (autostability spectrum) of a computable structure  $\mathfrak{M}$  is the set

$$\text{CatSpec}(\mathfrak{M}) = \{ \mathbf{d} : \mathfrak{M} \text{ is } \mathbf{d}\text{-computably categorical} \}.$$

A Turing degree  $\mathbf{d}_0$  is the **degree of categoricity** of  $\mathfrak{M}$  if  $\mathbf{d}_0$  is the least degree in  $\text{CatSpec}(\mathfrak{M})$ .

# Spectrum of decidable categoricity

Definition(Goncharov, 2011)

The **decidable categoricity spectrum** (autostability spectrum relative to strong constructivizations) of the structure  $\mathfrak{M}$  is the set

$$\text{DecCatSpec}(\mathfrak{M}) = \{\mathbf{d} : \mathfrak{M} \text{ is decidable } \mathbf{d}\text{-categorical}\}.$$

A Turing degree  $\mathbf{d}_0$  is the **degree of decidable categoricity** of  $\mathfrak{M}$  if  $\mathbf{d}_0$  is the least degree in  $\text{DecCatSpec}(\mathfrak{M})$ .

## Prime models and complete formulas

Let  $\mathfrak{M}$  be a structure of a signature  $\sigma$ .  $\text{Th}(\mathfrak{M})$  denotes the first-order theory of  $\mathfrak{M}$ .

A structure  $\mathfrak{M}$  is a **prime model** (of the theory  $\text{Th}(\mathfrak{M})$ ) if  $\mathfrak{M}$  is elementary embeddable into every model  $\mathfrak{N}$  of the theory  $\text{Th}(\mathfrak{M})$ .

A structure  $\mathfrak{M}$  is an **almost prime model** if there exists a finite tuple  $\bar{c}$  from  $\mathfrak{M}$  such that  $(\mathfrak{M}, \bar{c})$  is a prime model.

A first-order formula  $\psi(x_0, \dots, x_n)$  is a **complete formula** for the theory  $\text{Th}(\mathfrak{M})$  if  $\mathfrak{M} \models \exists \bar{x} \psi(\bar{x})$  and, for every  $\sigma$ -formula  $\varphi(\bar{x})$ , either  $\mathfrak{M} \models \forall \bar{x} (\psi(\bar{x}) \rightarrow \varphi(\bar{x}))$  or  $\mathfrak{M} \models \forall \bar{x} (\psi(\bar{x}) \rightarrow \neg \varphi(\bar{x}))$ .

# Nurtazin's criterion

## Theorem (Nurtazin 1974)

Suppose that  $\mathfrak{M}$  is a decidable structure of a signature  $\sigma$ .  $\mathfrak{M}$  is decidable categorical if and only if there exists a finite tuple  $\bar{c}$  from  $\mathfrak{M}$  such that the following holds:

- (a) The structure  $(\mathfrak{M}, \bar{c})$  is a prime model of the theory  $\text{Th}(\mathfrak{M}, \bar{c})$ .
- (b) Given a  $(\sigma \cup \{\bar{c}\})$ -formula  $\psi(\bar{x})$  one can effectively, uniformly in  $\psi$ , determine whether  $\psi$  is a complete formula for  $\text{Th}(\mathfrak{M}, \bar{c})$ .

# Goncharov's result

## Theorem (Goncharov, 2011)

Let  $\mathbf{d}$  be a Turing degree. Suppose that  $\mathfrak{M}$  is a decidable structure of a language  $L$ ,  $\bar{a}$  is a finite tuple from  $\mathfrak{M}$  such that the following conditions hold.

- (a) The structure  $(\mathfrak{M}, \bar{a})$  is a prime model.
- (b) Given a  $(L \cup \{\bar{a}\})$ -formula  $\psi(\bar{x})$ , one can effectively relative to  $\mathbf{d}$ , uniformly in  $\psi$ , determine whether  $\psi$  is a complete formula in the theory  $Th(\mathfrak{M}, \bar{a})$ .

Then  $\mathfrak{M}$  is decidable  $\mathbf{d}$ -categorical.

# Known results

- ▶ [Goncharov \(2011\)](#) Every c.e. degree  $\mathbf{d}$  is the degree of decidable categoricity of some decidable almost prime model of infinite signature.
- ▶ [Bazhenov](#)
  - ▶ For every computable ordinal  $\alpha$ , the Turing degree  $0^{(\alpha)}$  is a degree of decidable categoricity for some decidable Boolean algebra. [\(2016\)](#)
  - ▶ For a computable successor ordinal  $\alpha$ , every Turing degree c.e. in and above  $0^{(\alpha)}$  is the degree of decidable categoricity for some decidable structure. [\(2016\)](#)
  - ▶ For an infinite computable successor ordinal  $\beta$ , every Turing degree c.e. in and above  $0^{(\beta)}$  is the degree of decidable categoricity for some linear order. [\(2017\)](#)
  - ▶ The set of all PA-degrees is the decidable categoricity spectrum. [\(2016\)](#).

Let  $T$  be a complete theory of a signature  $\sigma$  and  $p$  be the set of  $\sigma$ -formulas in free variables  $x_1, \dots, x_n$ . We call  $p$  an (complete) **n-type** if the following holds:

- (a)  $p \cup T$  is satisfiable
- (b)  $\varphi \in p$  or  $\neg\varphi \in p$  for all  $\sigma$ -formulas in free variables  $x_1, \dots, x_n$ .

We let  $\mathbf{S}_n(\mathbf{T})$  be the set of all  $n$ -types of  $T$ .

$n$ -type  $p$  is said to be **principal** if it contains a complete formula.

# Stone topology

Let  $T$  be a complete theory of a signature  $\sigma$ . For  $\sigma$ -formula in free variables  $x_1, \dots, x_n$ , which is satisfiable with  $T$ , let

$$[\varphi] = \{p \in S_n(T) : \varphi \in p\}$$

The **Stone topology** (Mal'cev topology) on  $S_n(T)$  is generating by taking the sets  $[\varphi]$  as basic open sets.

## Cantor-Bendixson rank

For a topological space  $X$  and an ordinal  $\alpha$ , the  $\alpha$ -th **Cantor-Bendixson derivative** of  $X$  is defined by transfinite induction as follows, where  $X'$  is the set of all limit points of  $X$ :

- ▶  $X^0 = X$
- ▶  $X^{\alpha+1} = (X^\alpha)'$
- ▶  $X^\lambda = \bigcap_{\alpha < \lambda} X^\alpha$ , for limit ordinals  $\lambda$ .

The smallest ordinal  $\alpha$ , such that  $X^{\alpha+1} = X^\alpha$ , is called the **Cantor-Bendixson rank** (CB-rank) of  $X$ .

Given a type  $p \in S_n(T)$ , we say that its **CB-rank** is  $\alpha$ , written  $CB(p) = \alpha$ , if  $p \in X^\alpha \setminus X^{\alpha+1}$ .

Equivalently,  $CB(p) = \alpha \Leftrightarrow p$  is an isolated point of  $(S_n(T))^\alpha$ .  
 $CB(p) = 0 \Leftrightarrow p$  is a principal type of  $T$ .

# $\Sigma_2^0$ Turing degree

## Theorem 1

For every  $\Sigma_2^0$  Turing degree  $\mathbf{d}$ , such that  $\mathbf{d} \geq 0'$  there exists a decidable model  $\mathfrak{M}$  such that:

- (a) the set of complete formulas of  $Th(\mathfrak{M})$  is computable,
- (b)  $\mathfrak{M}$  is not homogeneous and realizes 1-types only of Cantor-Bendixson rank  $\leq 1$
- (c)  $\mathbf{d}$  is a degree of decidable categoricity of  $\mathfrak{M}$  .

## Theorem 2 (with N. Bazhenov)

Let  $\mathfrak{M}$  be a model that realizes types only of Cantor-Bendixon rank  $\leq 1$  then  $\mathfrak{M}$  is almost prime.

Thank you for your attention!