
LOG | COLLOQUIUM

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Tree-shaped one-pass tableau systems for Linear Temporal Logic satisfiability checking

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Joint work with Angelo Montanari, Mark Reynolds, Luca Geatti

The need for formal verification

Safety-critical systems need to avoid bugs at all costs.

Formal Verification develops automatic techniques to provide mathematical proofs of software correctness.

Formal verification and Logic

In formal verification, an **abstract model** of the system is checked against a **formal specification** of the desired behavior.

- Systems are usually modeled as **automata**.
- Specifications are usually expressed as **temporal logic** formulas.

This is the **model checking** problem.

Linear Temporal Logic

Linear Temporal Logic (LTL) is a propositional modal logic commonly used as specification language.

$X\alpha$	α will be true at the next state.
$\alpha U \beta$	β will eventually be true, and α always holds until then.
$F\beta \equiv T U \beta$	β will eventually be true.
$G\beta \equiv \neg F \neg \beta$	β will always be true.

Linear Temporal Logic (2)

*If infinitely many requests are received,
then infinitely many replies are sent.*

$$GF r \longrightarrow GF q$$

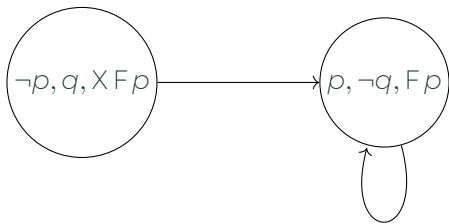
LTL satisfiability checking

LTL **satisfiability** is the problem of checking whether there exists a model that satisfies a given LTL formula.

- Important to check **consistency** of specifications.
- Seemingly very hard problem: **PSPACE**-complete.
- Many algorithmic solutions to solve it:
 - (Büchi) Automata-based
 - **Tableau** methods
 - Temporal resolution
 - Reduction to model checking
 - ...

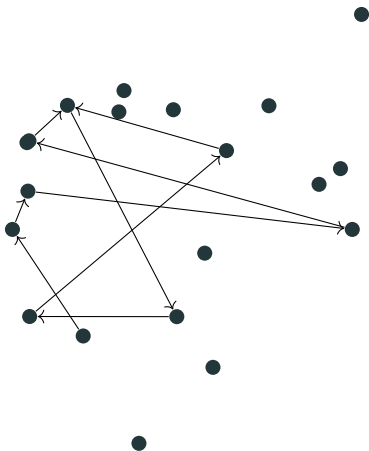
Tableaux-based methods for LTL satisfiability

Common tableaux-based methods for LTL are **graph-shaped** and require multiple passes to be built and checked for successful paths.



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A One-Pass Tree-Shaped Tableau for LTL

A **one-pass tree-shaped** tableau method for LTL satisfiability was recently proposed.

GandALF 2016

M. Reynolds. "A New Rule for LTL Tableaux." In: *Proc. of the 7th International Symposium on Games, Automata, Logics and Formal Verification*. GandALF 2016

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M. Bertello, N. Gigante, A. Montanari, and M. Reynolds. "Leviathan: A New LTL Satisfiability Checking Tool Based on a One-Pass Tree-Shaped Tableau." In: *Proc. of the 25th International Joint Conference on Artificial Intelligence*. IJCAI 2016

<http://www.github.com/corralx/leviathan>

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A One-Pass Tree-Shaped Tableau for LTL

A **one-pass tree-shaped** tableau method for LTL satisfiability was recently proposed

- Purely tree-shaped rule-based search procedure.
- A single pass is sufficient to determine the acceptance or rejection of a given branch.
- Very **simple** set of rules, combining the simplicity of declarative tableaux with the efficiency of one-pass systems.
- Easy to **extend**.
- Easy to parallelize with huge speedups.

Example

$\{G F(p \wedge X \neg p)\}$

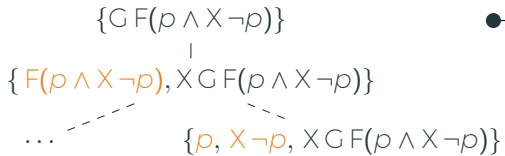


Example

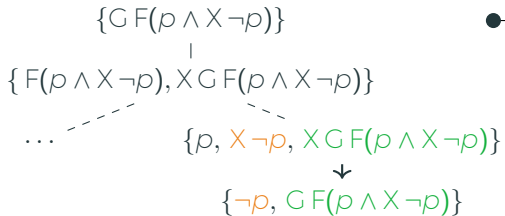
$$\begin{array}{c} \{G F(p \wedge X \neg p)\} \\ | \\ \{F(p \wedge X \neg p), X G F(p \wedge X \neg p)\} \end{array}$$



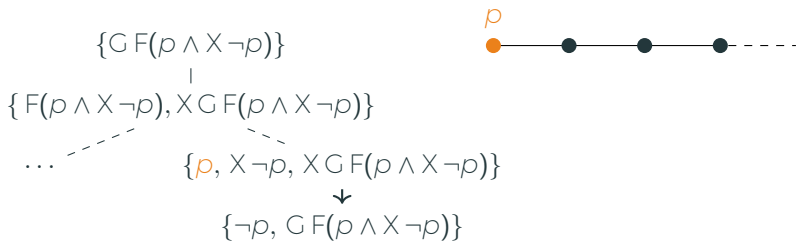
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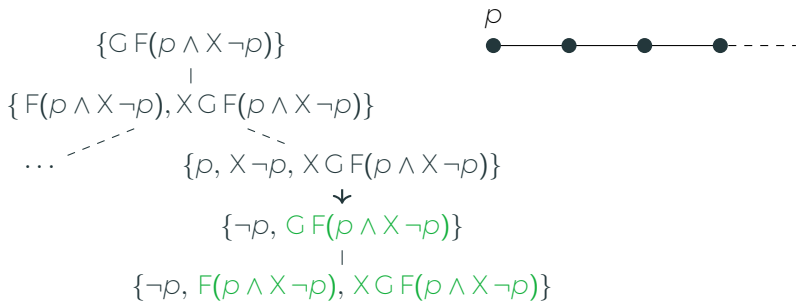
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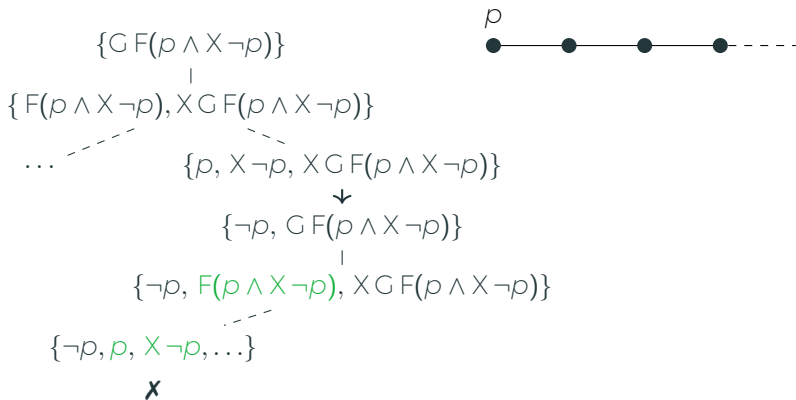
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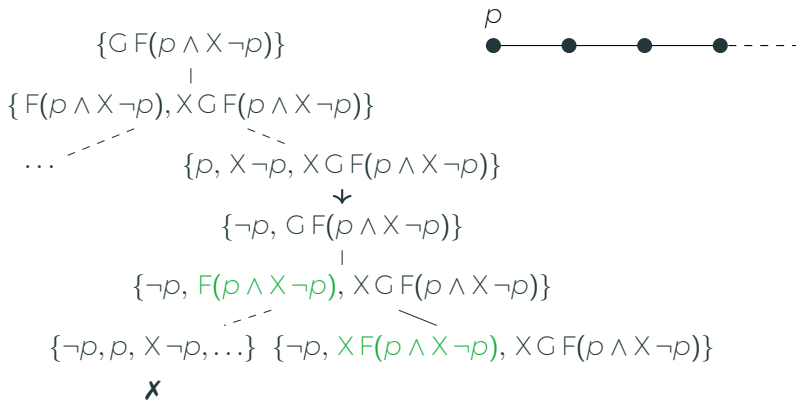
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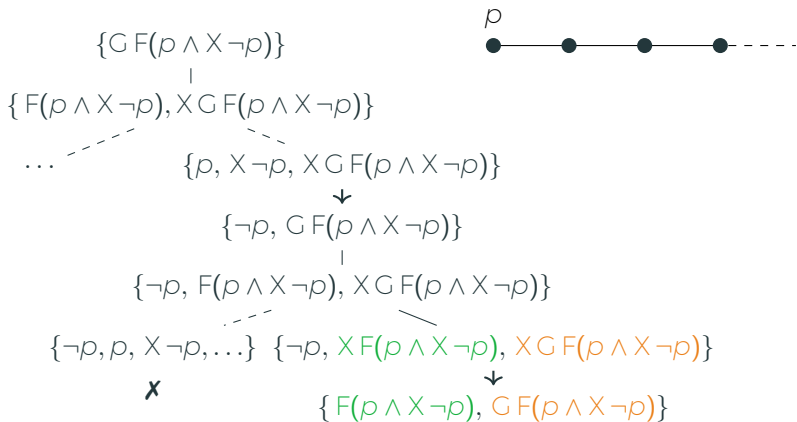
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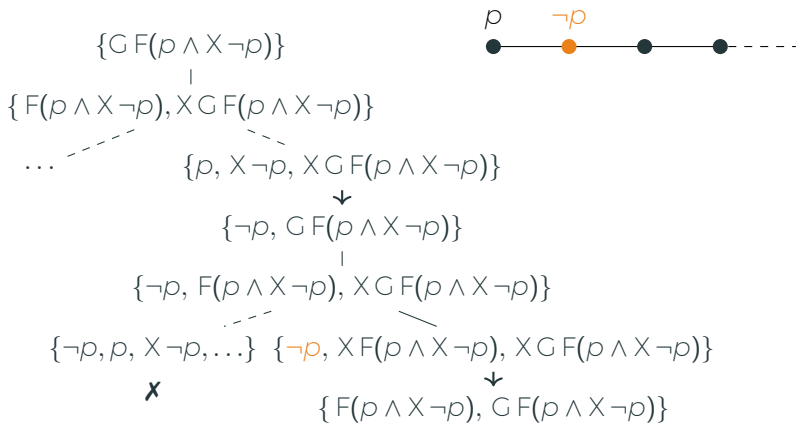
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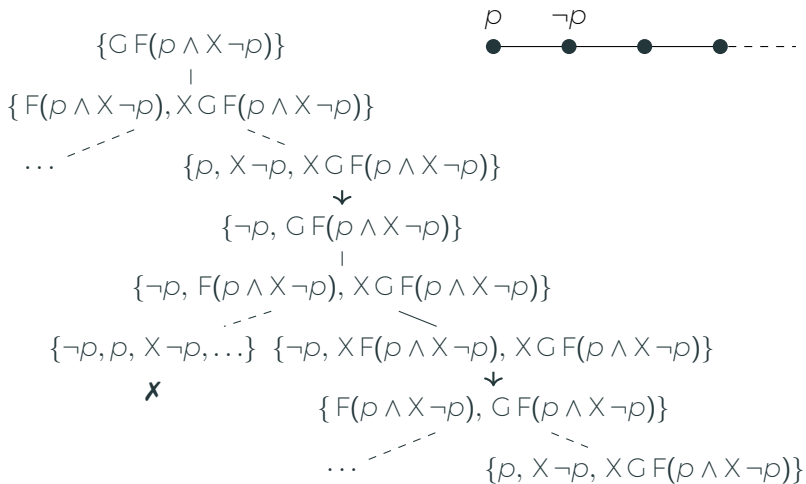
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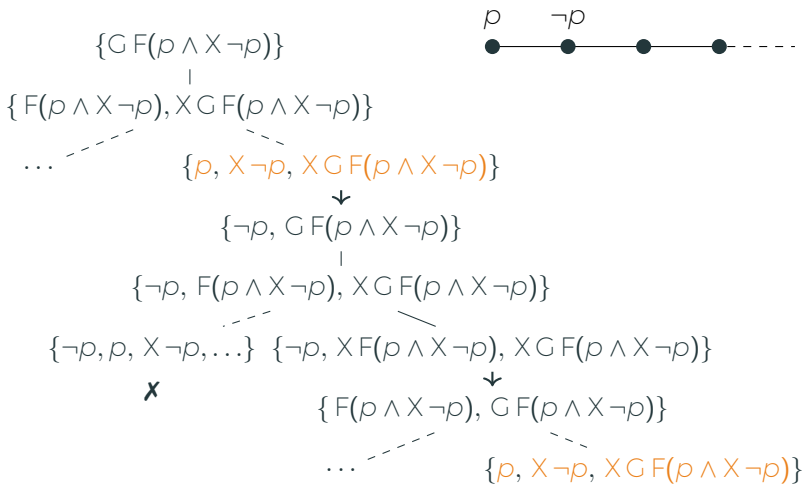
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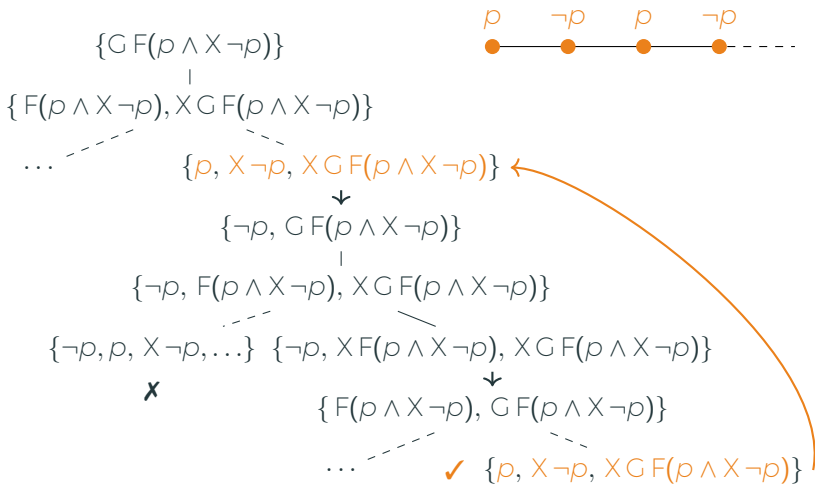
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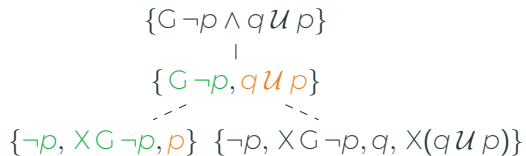
Example - unsatisfiable formula

$$\{G \neg p \wedge q \mathcal{U} p\}$$

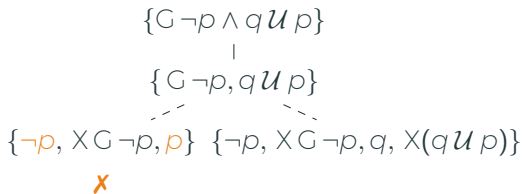
Example - unsatisfiable formula

$$\begin{array}{c} \{G \neg p \wedge q \mathcal{U} p\} \\ | \\ \{G \neg p, q \mathcal{U} p\} \end{array}$$

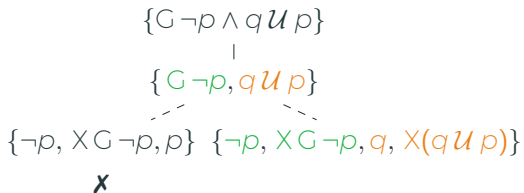
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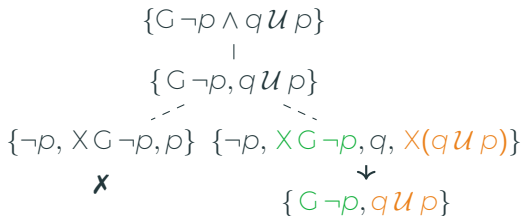
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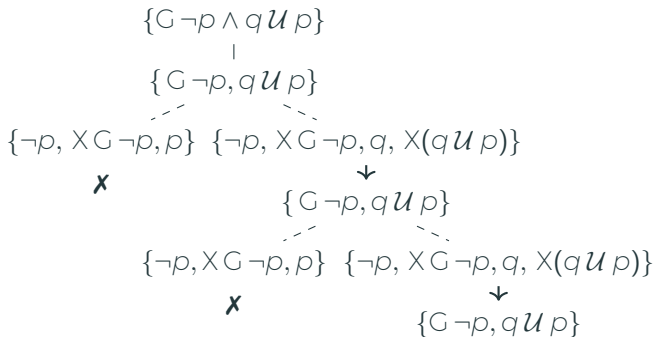
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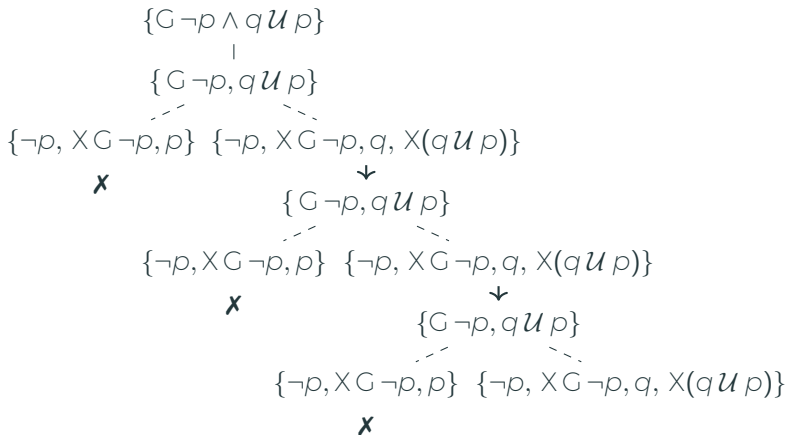
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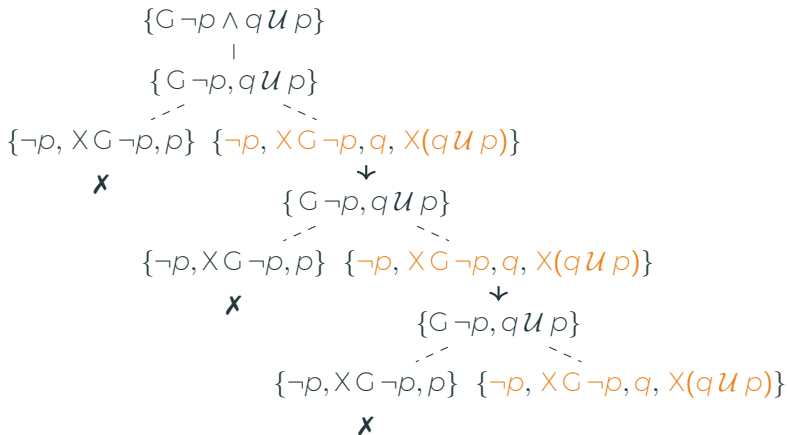
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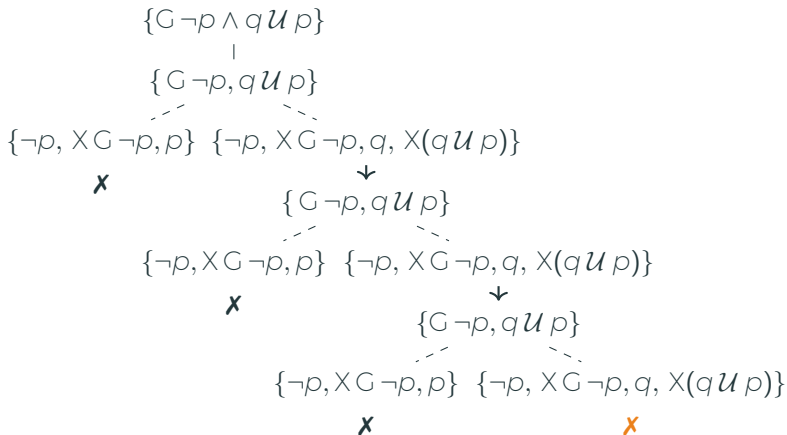
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Beyond LTL

Our conjecture is that the simplicity of the tableau can be adapted to more expressive logics.

Extensions developed so far:

- LTL with **past operators**;

LPAR 21

Nicola Gigante, Angelo Montanari, and Mark Reynolds. "A One-Pass Tree-Shaped Tableau for LTL+Past." In: *21st Int. Conference on Logic for Programming, Artificial Intelligence and Reasoning*. 2017, pp. 456-473

Beyond LTL

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Extensions developed so far:

- LTL with **past operators**;
- Timed Propositional Temporal Logic (TPTL) and Bounded TPTL + Past ($TPTL_b+P$)

Submitted

Luca Geatti, Nicola Gigante, Angelo Montanari, and Mark Reynolds. "One-pass and tree-shaped tableau systems for TPTL and $TPTL_b+P$." In: *Submitted for review*. 2018

Beyond LTL

Our conjecture is that the simplicity of the tableau can be adapted to more expressive logics.

Future and in progress work:

- extension to full TPTL with past;
- translation of tableau rules into new SAT/SMT encodings for the LTL satisfiability problem;
- study of the model-theoretic properties of models found by the tableau PRUNE rule.

Future work

Future lines of work:

- Add the past to our satisfiability checking tool.
 - Not trivial: our rule causes a lot of backtracking
- Exploit the modular structure of the tableau to extend it to other LTL extensions.
- Implement these extensions: one tool for a broad family of linear time logics.

Thank you!

Questions?

Bibliography I

- [Ber+16] M. Bertello, N. Gigante, A. Montanari, and M. Reynolds. “Leviathan: A New LTL Satisfiability Checking Tool Based on a One-Pass Tree-Shaped Tableau.” In: *Proc. of the 25th International Joint Conference on Artificial Intelligence*. IJCAI 2016.
- [Gea+18] Luca Geatti, Nicola Gigante, Angelo Montanari, and Mark Reynolds. “One-pass and tree-shaped tableau systems for TPTL and TPTL_b+P.” In: *Submitted for review*. 2018.
- [GMR17] Nicola Gigante, Angelo Montanari, and Mark Reynolds. “A One-Pass Tree-Shaped Tableau for LTL+Past.” In: *21st Int. Conference on Logic for Programming, Artificial Intelligence and Reasoning*. 2017, pp. 456–473.
- [MR17] John Christopher McCabe-Dansted and Mark Reynolds. “A Parallel Linear Temporal Logic Tableau.” In: *Proceedings 8th International Symposium on Games, Automata, Logics and Formal Verification*. 2017, pp. 166–179.
- [Rey16] M. Reynolds. “A New Rule for LTL Tableaux.” In: *Proc. of the 7th International Symposium on Games, Automata, Logics and Formal Verification*. GandALF 2016.

Why three occurrences?

Consider this formula:

$$\begin{aligned}\phi \equiv & p \wedge G(p \leftrightarrow X\neg p) \wedge GFq_1 \wedge GFq_2 \wedge \\ & G\neg(q_1 \wedge q_2) \wedge G(q_1 \rightarrow \neg p) \wedge G(q_2 \rightarrow \neg p)\end{aligned}$$

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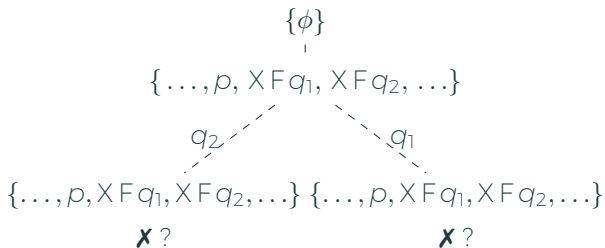
$$\begin{array}{c}\{\phi\} \\ | \\ \{\dots, p, XFq_1, XFq_2, \dots\}\end{array}$$

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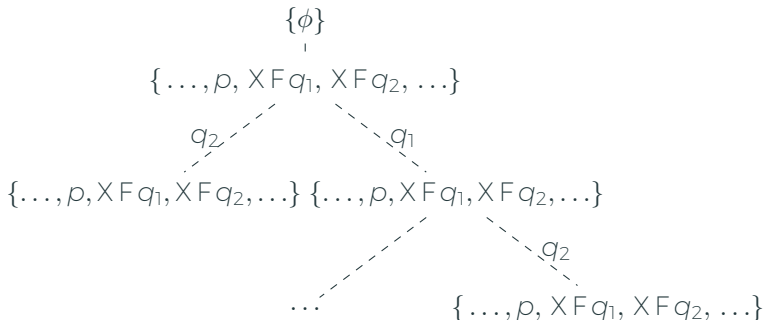


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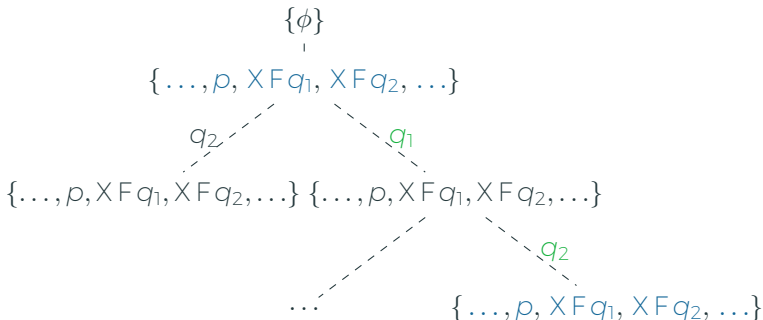


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